

# The `semantic` package<sup>†</sup>

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## Abstract

The aim of this package is to help people doing programming languages using  $\LaTeX$ . The package provides commands that facilitates the use of the notation of semantics and compilation in your documents. It provides an easy way to define new ligatures, *eg* making `=>` a short hand for `\RightArrow`. It facilitates the drawing of inference rules and allows you to draw T-diagrams in the `picture` environment. It supports writing extracts of computer languages in a uniform way. It comes with a predefined set of shorthand suiting most people.

## Contents

<b>1</b>	<b>Loading</b>	<b>2</b>	<b>3.2</b>	<b>Formatting the Entries . . .</b>	<b>4</b>
<b>2</b>	<b>Math Ligatures</b>	<b>2</b>	<b>4</b>	<b>T-diagrams</b>	<b>5</b>
2.1	Defining New Math Ligatures . . . . .	2	<b>5</b>	<b>Reserved Words</b>	<b>7</b>
2.2	Turning Math Ligatures On and Off . . . . .	3	5.1	Bells and Whistles: Spacing in Math Mode . .	7
2.3	Protecting Fragile Math Commands . . . . .	3	<b>6</b>	<b>Often Needed Short Hands</b>	<b>8</b>
<b>3</b>	<b>Inference Rules</b>	<b>3</b>	6.1	The Meaning of: <code>[[ and ]]</code> .	8
3.1	Controlling the Appearance . . . . .	4	6.2	Often Needed Symbols . .	8
			<b>7</b>	<b>Some Notes about the Files</b>	<b>9</b>

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`semantic` is a  $\LaTeX 2_{\epsilon}$  package facilitating the writing of notation of programming languages and compilation. To use it, the file `semantic.sty` should

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\*This file has version 2.0 $\epsilon$  and is dated 2003/10/28. It is CVS revision 1.11, dated 2003/10/28 13:45:57.

<sup>†</sup>Michael John Downes of AMS provided a patch to make the `semantic` compatible with `amsmath` v.2.01.

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be placed so that L<sup>A</sup>T<sub>E</sub>X can find it.

The `semantic` package consists of several parts, which can be used independently. The different parts are

**Ligatures** providing an easy way to define ligatures for often used symbols like  $\Rightarrow$  and  $\vdash$ .

**Inference Rules** facilitating the presentation of inference rules and derivations using inference rules.

**T-diagrams** providing T-diagrams as an extension the `picture` environment.

**Reserved word**<sup>1</sup> facilitating getting a uniform appearance of language constructions.

**Short hands** for often used symbols.

In the following we describe the use of the various parts of `semantic` and the installation. We also give a short introduction to the two files `semantic.dtx` and `semantic.ins`.

This package is—like most other computer-programs—provided with several bugs, insufficiencies and inconsistencies. They should be regarded as features of the package. To increase the excitement of using the package these features appear in unpredictable places. If they however get too annoying and seriously reduce your satisfaction with `semantic`, please notify us. You could also drop us a note if you would like to be informed when `semantic` is updated.

## 1 Loading

There is two ways of loading the `semantic` package. You can either load it with all the parts, or to save time and space, you can load, only the parts you will use.

In the first case you just include

```
\usepackage{semantic}
```

in your document preamble.

In the other case you include

```
\sepackage[parts]{semantic}
```

in your document preamble. `parts` is a comma separated list of the parts you wants to include. The possibilities are: `ligature`, `inference`, `tdiagram`, `reserved`, and `shorthand`. The different parts are described in detail below.

## 2 Math Ligatures

### 2.1 Defining New Math Ligatures

`\mathlig` When the package is loaded, you can define new ligatures for use in the math environments by using the `\mathlig{character sequence}{ligature commands}` command. `character sequence` is a sequence of characters<sup>2</sup> that must be entered in the source file to achieve the effect of the `ligature command`. If for example you write `\mathlig{-><-}{\rightarrow\leftarrow}`, subsequently typing `-$-><-$` will produce  $\rightarrow\leftarrow$ .

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<sup>2</sup>There are some restrictions on the characters you can use. This should be described here but isn't; basically you should stick to using the characters ' " ' ~ ! ? @ \* ( ) [ ] < > - + = | : ; . , / 0 . . . 9, and certainly this should suffice for any sane person.

## 2.2 Turning Math Ligatures On and Off

`\mathligson` By default, math ligatures are in effect when the `mathlgs` package is loaded, but this can be turned off and on by using the commands `\mathligsoff` and `\mathligson`. Thus, typing `'$-><-$ \mathligsoff $-><-$ \mathligson $-><-$'` will produce  $\rightarrow\leftarrow - > < - \rightarrow\leftarrow$ .

## 2.3 Protecting Fragile Math Commands

`\mathligprotect` Unfortunately, some macros used in math mode will break when using `mathlgs`, so they need to be turned into protected macros with the declaration `\mathligprotect{<macro>}`. *NOTE:* This declaration only needs to be issued once, best in the preamble.

## 3 Inference Rules

`\inference` Inference rules like  
`\inference*`

$$\text{It(1): } \frac{\rho \vdash E \Rightarrow \text{FALSE}}{\rho \vdash \text{while } E \text{ do } s \Rightarrow \rho} \quad \text{It(2): } \frac{\rho \vdash E \Rightarrow \text{TRUE} \quad \rho \vdash s \Rightarrow \rho'}{\rho \vdash \text{while } E \text{ do } s \Rightarrow \rho''}$$

and

$$\rightarrow^*_1 \frac{\begin{array}{l} p, M \rightarrow^* p', M' \\ p', M' \rightarrow p'', M'' \end{array}}{p, M \rightarrow^* p'', M''} \quad \rightarrow^*_2 \frac{}{p, M \rightarrow^* p, M}$$

are easily set using `\inference` and `\inference*`. The syntax is

`\inference[<name>]{<line1> \ \cmd{\littdots} \ \ <linen>}{<conclusion>}`

and

`\inference*[<name>]{<line1> \ \cmd{\littdots} \ \ <linen>}{<conclusion>}`

where  $n \geq 0$  so that you can also type axioms. When using `\inference` the bar will be as wide as the conclusion and the premise, whichever is widest; while `\inference*` only will make the bar as wide as the conclusion (It(2) above). The optional names are typeset on the side of the inferences that they appear.

Each line consists of premises separated by `&`:

`<premise1>&\cmd{\littdots}&<premisem>`

Note that  $m$  can also be zero, which is used when typing axioms. Each premise and the conclusion are by default set in math mode (*see* however 4).

The rules are set so that the line flushes with the center of small letters in the surrounding text. In this way, secondary conditions or names (like the first example above) can be written in the surrounding text. One may also set the rules in a table as shown below:

$$\begin{array}{l} \text{Transitive (1): } \\ \text{Transitive (2): } \end{array} \frac{\begin{array}{l} p, M \rightarrow^* p', M' \\ p', M' \rightarrow p'', M'' \end{array}}{p, M \rightarrow^* p'', M''} \quad \frac{}{p, M \rightarrow^* p, M}$$

An inference rule can be nested within another rule without problems, like in:

$$\frac{\frac{\frac{\rightarrow^*_2 \overline{p, M \rightarrow^* p, M}}{p, M \rightarrow^* p', M'}}{\rightarrow^*_1} \quad p', M' \rightarrow p'', M''}{\rightarrow^*_1 \overline{p, M \rightarrow^* p', M'}}}{\rightarrow^*_1 \overline{p, M \rightarrow^* p'', M''}}$$

### 3.1 Controlling the Appearance

`\setpremissesend`  
`\setpremissesspace`  
`\setnamespace`

The appearance of the inferences rules can be partly controlled by the following lengths:

$$\frac{\underbrace{\text{namespace}}_{\text{name}} \overbrace{\text{premissesend} \quad \text{premissesspace}}^{\text{premise} \quad \text{premise}}}{\text{conclusion}}$$

The lengths are changed using the three commands `\setnamespace{⟨length⟩}`, `\setpremissesend{⟨length⟩}` and `\setpremissesspace{⟨length⟩}`. `⟨length⟩` can be given in both absolute units like `pt` and `cm` and in relative units like `em` and `ex`. The default values are:  $1\frac{1}{2}\text{em}$  for `premissesspace`,  $\frac{3}{4}\text{em}$  for `premissesend` and  $\frac{1}{2}\text{em}$  for `namespace`. Note that the lengths *cannot* be altered using the ordinary L<sup>A</sup>T<sub>E</sub>X-commands `\setlength` and `\addtolength`.

Besides that, the appearance of inference rules is like fractions in math: Among other things the premisses will *normally* be at same height above the baseline and there is a minimum distance from the line to the bottom of the premisses.

⚡ Fetching the font information from the math font and the evaluation (in case they are defined in relative units) of the lengths mentioned above is done just before the individual rule is set. This is demonstrated by the following construction (which admittedly is not very useful):

$$\frac{\text{Large} \overbrace{\text{normalsize} \overbrace{\text{footnotesize} \overbrace{\text{tiny} \overline{\text{Conclusion}}}}^{\text{Conclusion}}}}^{\text{Conclusion}}}{\text{Conclusion}}$$

Note that from top to bottom, the leaves get bigger and the names get further from the line below.

### 3.2 Formatting the Entries

`\predicate`

⚡ To set up a single predicate (a premise or conclusion) the single-argument command `\predicate` is used. This allows a finer control of the formatting. As an example, all premisses and conclusions can be set in mathematics mode by the command:

`\renewcommand{\predicate}[1]{\$ #1 \$}`

⚡ `semantic` uses `\predicate` on a premise only when the premise does not contain a nested `\inference`.<sup>3</sup> So even if the declaration above has been given, `\inference` is *never* be executed in math mode. Neither is it used on the premisses if you write:

`\inference{\inference...}\{...\}`

`\predicatebegin`  
`\predicateend`

⚡ The default definition of `\predicate` is `\predicatebegin #1\predicateend`, where `\predicatebegin` and `\predicateend` are defined to ‘\$’. In this way the premisses and conclusions are set in math

<sup>3</sup>What `semantic` precisely does is to append the tokens `\inference \end` to the code of a premise, when it has isolated it. `semantic` then uses T<sub>E</sub>X’s pattern matching to search this new list of tokens for an appearance of the token `\inference`. When this is found the following token is examined, and if it is `\end`, `semantic` concludes that the premise does not contain a nested inference rule

$\diamond$  The motivation for introducing `\predicatebegin` and `\predicateend` was, however, to use TeX's pattern matching on macro arguments to do even more sophisticated formatting by redefining `\predicatebegin`. If for example, *every expression* is to be evaluated in an *environment* giving a *value*, and you would like to set *all the environment's values* in mathematics and the *expressions* in `typewriter`-font, then this could be facilitated by the definition:

```

\def\predicatebegin#1|-#2=>#3#4\predicateend{%
  #1 \vdash\texttt{#2}\stackrel{#3}{\rightarrow}_S #4$}

```

Then the inference (borrowed from M. HENNESSY, *The Semantics of Programming Languages*)

$$\text{TIR} \frac{D \vdash s \xrightarrow{v}_S s' \quad D \vdash s \xrightarrow{v'}_S s''}{D \vdash \text{Tl}(s) \xrightarrow{v'}_S s''}$$

can be accomplished by

```

\inference[TIR]{D |- $$s =>{v} s' & D |- $$s =>{v'} s''}
  {D |- Tl($$) =>{v'} s''}

```

Please note that the `ligatures` option *has not* been used above.

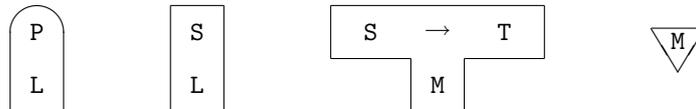
## 4 T-diagrams

```

\compiler
\interpreter
\program
\machine

```

To draw T-diagrams describing the result of using one or more compilers, interpreters etc., `semantic` has commands for the diagram:



These commands should only be used in a `picture` environment and are

```

\program{<program>,<implementation language>}
\interpreter{<source>,<implemenation language>}
\compiler{<source>,<machine>,<target>}
\machine{<machine>}

```

The arguments can be either a string describing the language (please do not begin the string with a macro name), or one of the four commands. However, combinations that make no sense—like implementing an interpreter on a program—are excluded, yielding an error message like:

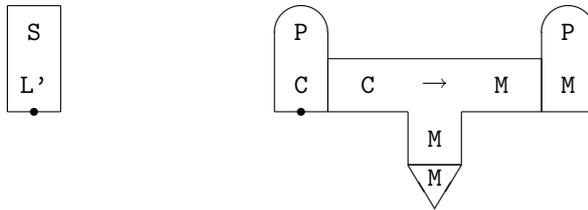
```

! Package semantic Error: A program cannot be at the bottom .

See the semantic package documentation for explanation.
Type H <return> for immediate help.
...

```

When you use a command as an argument `semantic`, will copy the language from the nested command and automatically place the two figures in proportion to each other. In this way, big T-diagrams can easily be drawn. The hole construction should be placed using a `\put` command, where the *reference point* is the center of the bottom of the figure corresponding to the outermost command. An example (with the reference point marked by  $\bullet$ ) will clarify some of these points. The figure



is obtained by the commands

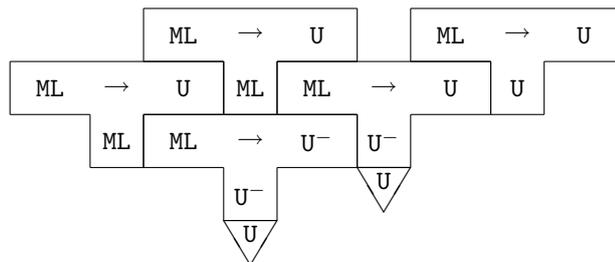
```
\begin{picture}(220,75)(0,-35)
\put(10,0){\interpreter{S,L'}}
\put(110,0){\program{P,\compiler{C,\machine{M},\program{P,M}}}}
\end{picture}
```

Note from the second example that when `\compiler` is used as “implementation language”-argument it is by convention attributed to the right of the figure. It is also worth mentioning that there is no strict demand on which command you should choose as the outermost, *ie* the second example could also be written (with a change of the parameters of `\put` due to the new reference point) as

```
\put(160,-20){\compiler{\program{P,C},\machine{M},\program{P,M}}}
```

starting off in the middle instead of using a “left-to-right”-approach. In fact, it is often easier to start in the middle, since this is where you get the least levels of nesting.

Even though most situations may be handled by means of nesting, it is in some rare cases adequate to use different language symbols on the two sides of the line of touch. When *eg* describing bootstrapping the poor U-code implementation can be symbolized by  $U^-$ , indicating that the poor implementation is still executed on a U-machine. This can be done by providing the symbol-command with an optional argument immediately after the command name. Thus the bootstrapping



is typed

```
\compiler{\compiler{ML,ML,U},\machine[U$^-]{U},\compiler{
\compiler{ML,ML,U},\machine[U$^-]{U},\compiler{ML,U,U}}}
```

For calculating the dimensions of the `picture`-environment, one needs the dimensions of the individual figures. In units of `\unitlength` they are the following:

compiler: 80*40	machine: 20*17.1
interpreter: 20*40	program: 20*40



will make **in** use the spacing of the relational symbols. The space command is applied to all the words in the reservation. Thus if you would like **in** and **let** to have different space commands, you must specify them in two different `\command`.

The drawback of using the math spacing is that in the rare cases where you use the reserved words in super- or subscripts, most of the spacing will disappear. This can be avoided by defining the replacement text to be the word plus a space, eg `\;in\;`. For this end a reservation of a word can be followed by an explicit replacement text in brackets, eg

```
\command{let[let\;], in[\;in\;]}
```

The formatting of `\command` (with the setting above: `\mathbf`) will still be used so it is only necessary to provide the replacement text. Note that each word in the reservation can have its own optional replacement text.

The drawback of this method is, that the you also get the space, if you use the reserved word “out of context”, for instance referring to the **in**-token! In these cases you can cancel the space by hand using `\!`.

This option is also usefull, if you want to typeset the same word in two different styles. If you for instance sometimes want ‘let’ to be typeset as a command and sometimes as data, you can define

```
\command{let}
\data{Let[Let]}
```

Then `\<let>` will typeset the word ‘let’ as a command, while `\<Let>` will typeset it as data. Note that in both cases the word appears in lower case.

Unfortunately there is no way to get the right spacing everytime, so you will have to choose which of the two methods serves you the best.

## 6 Often Needed Short Hands

Within the field of semantics there are a tradition for using some special. symbols. These are provided as default as short hand in the `semantic` package. Most of the following symbols are defined as ligatures, and hence the `ligature` option is always implied when the `shorthand` option is provided.

### 6.1 The Meaning of: `[[ and ]]`

`[[` The symbols for denoting the meaning of an expression, `[[ and ]]` are provided as short hands in math with the ligatures `|[ and ]|`.

### 6.2 Often Needed Symbols

The following ligatures are defined for often needed symbols

$\vdash$	$\dashv$	$\vDash$	$\Vdash$
$\longleftrightarrow$	$\leftrightarrow$	$\Leftrightarrow$	$\Leftrightarrow$
$\rightarrow$	$\rightarrow$	$\longrightarrow$	$\twoheadrightarrow$
$\Rightarrow$	$\Rightarrow$	$\implies$	$\implies$
$\leftarrow$	$\leftarrow$	$\longleftarrow$	$\twoheadleftarrow$
$\Leftarrow$	$\Leftarrow$	$\Leftarrow$	$\Leftarrow$

All the single directed arrows also comes in a starred and plussed form, eg `*<=>` gives `*<=>` and `-->+` gives `\longrightarrow^+`.

`\eval`      To support writing denotational, semantics the commands `\comp` and `\eval`  
`\comp` are provided to describe the evaluation of programs respectively expressions.  
They have the same syntax: `\comp{<command>}{<environment>}`, which yields  $\mathcal{C}[\![<command>]\!](<environment>)$ . If you need to describe more than one kind of evaluations, e.g. both  $\mathcal{E}$  and  $\mathcal{E}^*$ , you can provide an optional argument immediately after `\comp` or `\eval`, respectively. As an example a denotational rule for a sequencing two commands

$$\mathcal{C}[\![C1 ; C2]\!]d = d' \quad \text{if } \mathcal{C}[\![C1]\!]d = d'' \quad \text{and } \mathcal{C}[\![C2]\!]d'' = d'$$

can be typed

```
\[
\comp{C1 ; C2}{d} = \mathhtt{d'} \quad \quad
\texttt{if } \$\comp{C1}{d} = \mathhtt{d'}\texttt{' }$ and
\comp{C2}{d''} = \mathhtt{d'}\texttt{' }$
\]
```

`\evalsymbol`      As shown above, you can get the evaluation symbol in itself. This is done by  
`\compsymbol` or `\evalsymbol`, respectively. These commands can also be supplied  
with an optional argument, e.g. `\evalsymbol[*]` to get  $\mathcal{E}^*$ .

`\exe`      The result of executing a program on a machine with som data can be de-  
scribed using `\exe`, which has the syntax `\exe{<program>}[<machine>]{<data>}`.  
The third Futumara projection `cogen = \[[spec]](spec.spec)` can be written  
`\mathhtt{cogen} = \exe{spec}{spec.spec}`\$. As an alternative, you can also  
give the machine L explicit:

$$\mathcal{C}[\![\texttt{cogen}]\!] = \exe{spec}[L]{spec.spec}$$

This will result in: `cogen = \[[spec]]L(spec.spec)`

## 7 Some Notes about the Files

`semantic` is distributed in two files, `semantic.dtx` and `semantic.ins`. Of these two files, `semantic.dtx` is the most important, as it contains all the essentials—users guide, code and documentation of the code. `semantic.ins` is used only to guide `docstrip` in generating `semantic.sty` from `semantic.dtx`.

To get `\[` and `\]`, used in `\comp`, `\eval` and `\exe` `semantic`, tries to load the package `bbold` written by A. JEFFREY. If this is not installed on your system, the symbols are simulated by drawing together two sharps. However, we recommend that you get `bbold` from your nearest CTAN-archive.

In addition to the users guide, you can also get the fully documented code. You need this, however, if you want to see how the macros are implemented the macros or if you want to change some part of the package. You should start by editing `semantic.dtx` and remove the percentage signs from the four lines starting at Line 2794

```
% \AlsoImplementation
% \EnableCrossrefs
% \CodelineIndex
% \RecordChanges
```

After saving the changes, you should run  $\LaTeX$  twice on the edited file to get a correct table of contents. Then you generate the index and change history, using `makeindex`:

```
makeindex -s gind.ist semantic
makeindex -s gglo.ist -o semantic.gls semantic.glo
```

After another run of  $\LaTeX$ , then the documentation is ready for printing.

© At last the boring formal stuff: The package is protected by the The  $\LaTeX$  Project Public License (lpl). You are encouraged to copy, use, delete etc. the package (`semantic.dtx` and `semantic.ins`) as much as your heart, but if you modify the code (even locally), you should change the name to avoid confusion. Under all circumstances, the package is still: ©1995–2000 Peter Møller Neergaard and Arne John Glenstrup.